Unique Paper Code	:	32341502
Name of the Course	:	B.Sc. (H) Computer Science
Name of the Paper	:	Theory of Computation
Semester	:	V
Year of admission	:	2019 and onwards

Duration: Three Hours

Maximum Marks: 75

## **Instructions for Candidates:**

- i. Attempt any **FOUR** questions.
- ii. Each question carries equal marks.
- iii. Consider  $\Sigma = \{a b\}$  for all the questions unless specified otherwise.
- 1. Consider the language L, of all the words of length four or more having first two letters same as last two letters.

For the above language, perform the following:

- Write all the words of L with the length five or less
- Write the number of words having length six
- Construct the regular expression
- Build Finite Automaton (FA)
- 2. Prove that it is true for all the regular languages that complement of a regular language is also regular.

Construct the deterministic finite automaton (DFA) that recognizes the same language as the non-deterministic finite automaton (NFA) given below and also describe the language recognized by it.



Convert the following transition graph into its equivalent regular expression:



- 3. Consider the following languages:
  - $L_1$  = Language of all the words having '**b**' at second position
  - $L_2$  = Language of all the words having no two consecutive *a*'s

Construct Finite Automaton  $FA_1$  for  $L_1$ ,  $FA_2$  for  $L_2$ . Also construct regular expression and finite automata for the following:

- $L_1 + L_2$
- $L_1 \cap L_2$
- (L<sub>1</sub>)\*
- 4. For the language  $L_3: a^{n+m}b^mc^n$ ; where  $\Sigma = \{a \ b \ c\}$  and  $m, n \ge 1$ , using pumping lemma, prove that the language is not regular. For the above language, do the following:
  - Write a context free grammar (CFG) for L<sub>3</sub>, and construct parse tree for the word *aaabbc* using this CFG
  - Build a pushdown automaton (PDA) for L<sub>3</sub>
- 5. Consider the following context free grammars (CFGs):

 $G_{1}: S \rightarrow bS | aX$   $X \rightarrow bS | aY$   $Y \rightarrow aY | bY | a | b$   $G_{2}: S \rightarrow XaX | bX$   $X \rightarrow XaX | XbX | \Lambda$   $G_{3}: S \rightarrow A | AA$   $A \rightarrow B | BB$   $B \rightarrow abB | b | bb$   $G_{4}: S \rightarrow BABABA$   $A \rightarrow a | \Lambda$   $B \rightarrow b | \Lambda$ 

For the above CFGs, perform the following:

- Write a regular expression for the language represented by G<sub>1</sub>
- Convert G<sub>2</sub> into its equivalent CFG without null(**A**)-production
- Convert G<sub>3</sub> into its equivalent CFG without unit-production
- Convert G<sub>4</sub> into its equivalent Chomsky Normal Form (CNF)
- 6. Consider the language L<sub>4</sub>:  $a^n b^n c^n$  where  $\Sigma = \{a \ b \ c\}$  and  $n \ge 1$ , and perform the following:
  - Build a turing machine M<sub>1</sub>, that accepts L<sub>4</sub>
  - Build another turing machine M<sub>2</sub>, that accepts complement of L<sub>4</sub>
  - Is L<sub>4</sub> a recursive language or recursively enumerable language? Justify your answer
  - Is L<sub>4</sub> a context-free language? Justify your answer.